

# Verifying OS components

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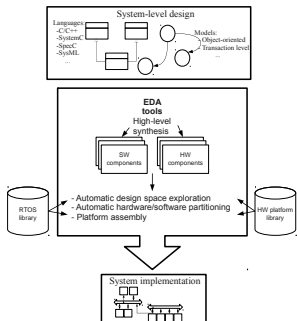
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Computer Architecture Group Presentations

## Why to Verify?

- Increasing the use of computational systems thus, increasing the need of them to be correct
  - “A system must do what it is *specified* to do”

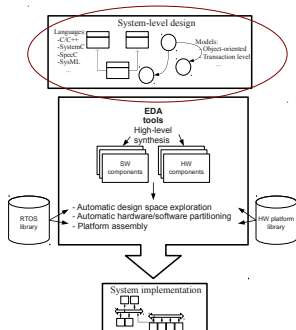
## What to Verify?



- Verification at System Level Design: regardless whether a component is gone to be implemented in SW or HW
- Chosen language: C++
  - HLS allows for “writing” HW in C++

Figure: Tomorrow's System-Level Design Flow

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Figure: Tomorrow's System-Level Design Flow

## How to Verify?

- Software Model Checking
- CBMC (C Bounded Model Checking)
  - Pointers
  - Arrays bound
  - User defined assertions

## Approach/Methodology

- Write the contract of the program/design to be verified
- Instrument the program with this contract
- Execute CBMC over the program to verify whether it follows its contract

## Writing the contract

- Hoare Logic:  $\{P\} C \{Q\}$
- Example (D's style)

```
void insert(T obj)
in
{
    size_at_pre = size ();
    assert(! ready.contains(obj));
}
out
{
    assert(size () == size_at_pre + 1);
    assert(ready.contains(obj));
}
body
{
    // Method's implementation
}
```

## Instrumenting the program with its contract

- Directly in the code
- Or using scenario adapters:

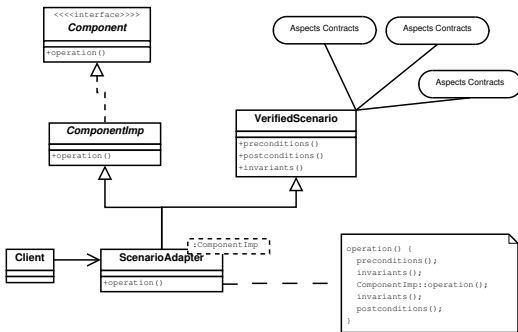


Figure: Verified Scenario



## Execute CBMC over the program to verify whether it follows its contract

- “Verification SUCCESSFUL”
- Or a counter-example, showing where the specification is not followed

# Case study: OS Scheduler

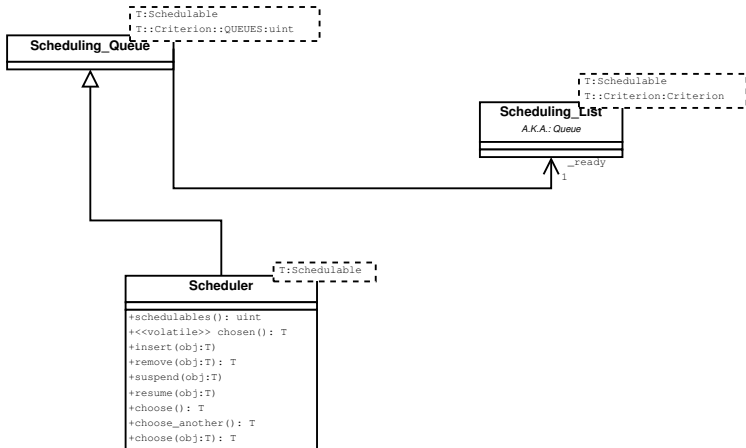


Figure: EPOS's Scheduler

# Scheduler's Contract

- Class invariants
  - **assert**(Base::size() >= 0)

## Scheduler's Contract

Method	Preconditions	Postconditions
<i>unsigned int schedulables()</i>	None	None
<i>T* volatile chosen()</i>	None	<b>assert</b> (ready.contains(var_chosen))
<i>void insert(T* obj)</i>	<b>assert</b> (! ready.contains(obj))	<b>assert</b> (Base::size() == size_at_pre + 1) <b>assert</b> (ready.contains(obj))
<i>T* remove(T* obj)</i>	None	<b>if</b> ready_at_pre.contains(obj) <b>then</b> <b>assert</b> (Base::size() == size_at_pre - 1) <b>else</b> <b>assert</b> (result = null) <b>endif</b> <b>assert</b> (! ready.contains(obj))
<i>void suspend(T* obj)</i>	None	<b>if</b> ready_at_pre.contains(obj) <b>then</b> <b>assert</b> (Base::size() == size_at_pre - 1) <b>endif</b> <b>assert</b> (! ready.contains(obj))
<i>void resume(T* obj)</i>	<i>same as insert</i>	<i>same as insert</i>
<i>T* choose()</i>	<b>assert</b> (ready.contains(var_chosen))	<b>assert</b> (ready.contains(var_chosen))
<i>T* choose_another()</i>	<b>assert</b> (ready.contains(var_another))	<b>assert</b> (ready.contains(var_another))
<i>T* choose(T* obj)</i>	<b>assert</b> (ready.contains(obj))	<b>assert</b> (ready.contains(var_chosen))